

Probabilistic Reasoning

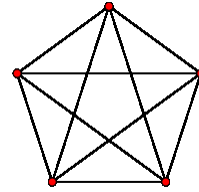
Unit # 11

How to Construct a JT

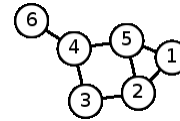
- Moralize the graph by adding links between parents of common children.
- Convert the given BN into an undirected graph by removing the arrowheads.
- Triangulate the graph.
 - If there is a cycle of length > 3 with no chord, add a chord between two non-adjacent vertices in the cycle.
- Order the nodes by maximum cardinality search.
 - Choose any node in the graph and label it 1.
 - For $i=2$ to n
 - Choose the node with the most labeled neighbors and label it i .
 - If any two labeled neighbors of i are not adjacent to each other
 - Add a link between those neighbors
 - Restart the algorithm.
- Find the cliques of the triangulated graph.
- Arrange as a junction tree.

Clique (From Wikipedia)

- In graph theory, a clique in an undirected graph $G = (V, E)$ is a subset of the vertex set $C \subseteq V$, such that for every two vertices in C , there exists an edge connecting the two.
- The size of a clique is the number of vertices it contains.
- Finding whether there is a clique of a given size in a graph (the clique problem) is NP-complete.

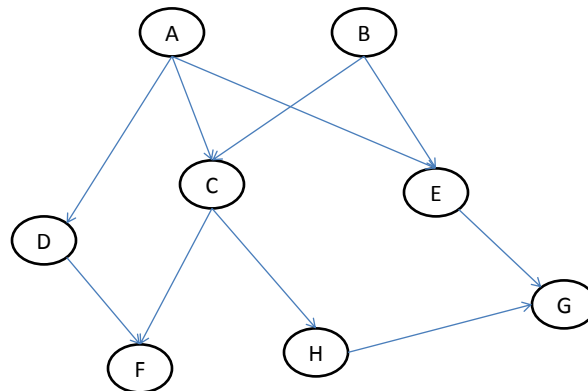


If a subgraph looks like this, the vertices in that subgraph form a clique of size 5.

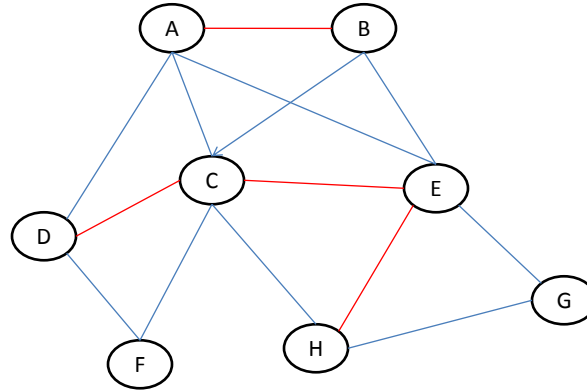


A graph with 6 vertices and 7 edges. On this graph vertices 1, 2, 5 is the only clique of size 3.

Junction Tree – Example I (Source: Laskey)



Junction Tree – Example I (Cont'd)

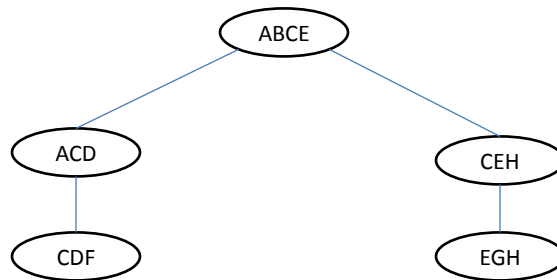


Sajjad Haider

Fall 2009

5

Junction Tree – Example I (Cont'd)

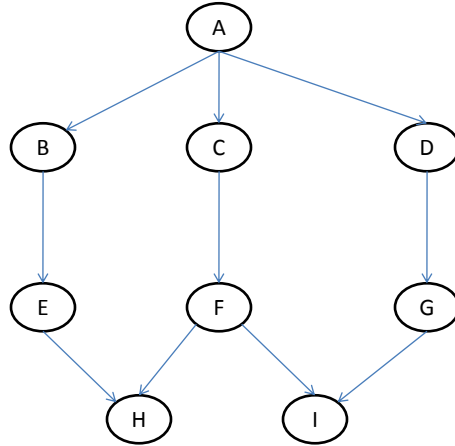


Sajjad Haider

Fall 2009

6

Junction Tree – Example II (Source: Korb & Nicholson)

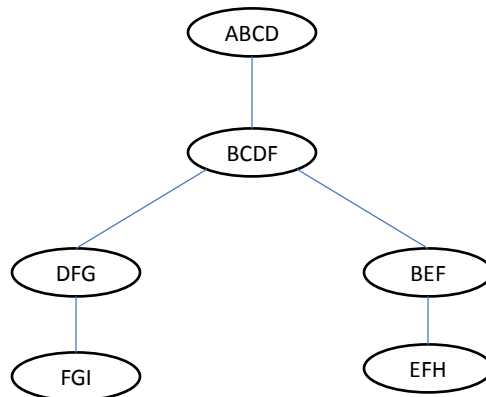


Sajjad Haider

Fall 2009

7

Junction Tree – Example II (Cont'd)



Sajjad Haider

Fall 2009

8

Approximate Inference Mechanisms

- Simulation Based Schemes
 - Logic Sampling
 - Likelihood Weighting
- Approximate Algorithms
 - Loopy Belief Propagation

Logic Sampling (Laskey)

- Logic sampling takes random draws from the joint distribution of all variables in the network
- Each node keeps a running count of how many times each of its states is drawn
- To take one random draw (one observation on all variables):
 - Sample the nodes in order (so parents of a node are sampled before the node)
 - For root node(s), select state with probability $P(\text{state})$
 - For non-root node(s) select state with probability $P(\text{state} | \text{sampled states of parents})$
 - After all nodes have been sampled
 - Check whether sampled values of evidence values match observed states
 - If they don't match, throw away the draw
 - If they match, increment the count for sampled state of each non-evidence node
- Estimate $P(X_t | X_e)$ by observed frequency among the retained samples

Algorithm (From Jensen and Nielsen)

1. Let (X_1, \dots, X_n) be a topological ordering of the variables.
2. For $j = 1$ to N :
 - a) For $i = 1$ to n :
 - Sample a state x_i for X_i using $P(X_i \mid \text{pa}(X_i) = \pi)$, where π is the configuration already sampled for $\text{pa}(X_i)$.
 - b) If $\mathbf{x} = (x_1, \dots, x_n)$ is consistent with \mathbf{e} , then

$$N(X_k = x_k) := N(X_k = x_k) + 1,$$

where x_k is the state that was sampled for X_k .

3. Return:

$$P(X_k = x_k \mid \mathbf{e}) \approx \frac{N(X_k = x_k)}{\sum_{x \in \text{sp}(X_k)} N(X_k = x)}.$$